CS 318 Principles of Operating Systems

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Pintos Virtual Memory Notes

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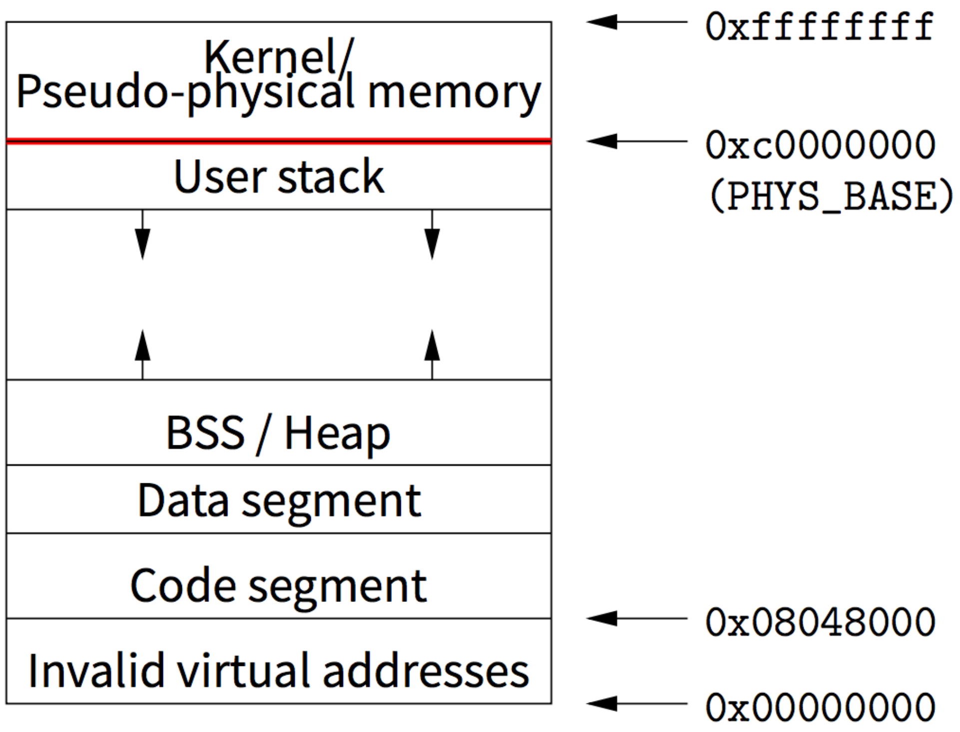
Pintos Virtual Memory Layout



i.e., the kernel lives in every process' address space!

A process' virtual address space is split into two regions

- The kernel lives in the high memory region, typically highest 1GB, i.e., from 3 to 4 GB.
- The user memory lives in the lower region, typically lower 3 GB, i.e., from 0 to 3 GB.

User Virtual Memory

Per process: a new page directory (pagedir) for each process

```
struct thread
  {
    tid_t tid;                          /* Thread identifier. */
    enum thread_status status;          /* Thread state. */
    char name[16];                      /* Name (for debugging purposes). */
    uint8_t *stack;                     /* Saved stack pointer. */
    int priority;                       /* Priority. */
    struct list_elem allelem;           /* List element for all threads list. */
    struct list_elem elem;              /* List element. */

#ifdef USERPROG
    /* Owned by userprog/process.c. */
    uint32_t *pagedir;                  /* Page directory. */
#endif

    /* Owned by thread.c. */
    unsigned magic;                     /* Detects stack overflow. */
  };
```

How Is A User Process Started?

```
$ pintos -p ../../examples/echo -a echo -- -f -q run 'echo cs318'
```

```
static void
run_task (char **argv)
{
  const char *task = argv[1];

  printf ("Executing '%s':\n", task);
#ifdef USERPROG
  process_wait (process_execute (task));
#else
  run_test (task);
#endif
  printf ("Execution of '%s' complete.\n", task);
}
```

How Is A User Process Started?

```
$ pintos -p ../../examples/echo -a echo -- -f -q run 'echo cs318'
```

```c
tid_t
process_execute (const char *file_name)
{
  char *fn_copy;
  tid_t tid;


  /* Make a copy of FILE_NAME.
     Otherwise there's a race between the caller and load(). */
  fn_copy = palloc_get_page (0);
  if (fn_copy == NULL)
    return TID_ERROR;
  strlcpy (fn_copy, file_name, PGSIZE);


  /* Create a new thread to execute FILE_NAME. */
  tid = thread_create (file_name, PRI_DEFAULT, start_process, fn_copy);
  if (tid == TID_ERROR)
    palloc_free_page (fn_copy);
  return tid;
}
```

"echo cs318"

Why?

The caller might free the file_name after this function returns!

,e.g., after you implement `exec`.

How Is A User Process Started?

```
$ pintos -p ../../examples/echo -a echo -- -f -q run 'echo cs318'
```

```
static void
start_process (void *file_name_)
{
  char *file_name = file_name_;
  struct intr_frame if_;
  bool success;

  /* Initialize interrupt frame and load executable. */
  memset (&if_, 0, sizeof if_);
  if_.gs = if_.fs = if_.es = if_.ds = if_.ss = SEL_UDSEG;
  if_.cs = SEL_UCSEG;
  if_.eflags = FLAG_IF | FLAG_MBS;
  success = load (file_name, &if_.eip, &if_.esp);

  /* If load failed, quit. */
  palloc_free_page (file_name);
  if (!success)
    thread_exit ();

  /* Start the user process by simulating a return from an interrupt */
  asm volatile ("movl %0, %%esp; jmp intr_exit" : : "g" (&if_) : "memory");
  NOT_REACHED ();
}
```

How Is A User Process Started?

```
$ pintos -p ../../examples/echo -a echo -- -f -q run 'echo cs318'
```

```
bool load (const char *file_name, void (**eip) (void), void **esp)
{
  struct thread *t = thread_current ();
  ...

  /* Allocate and activate page directory. */
  t->pagedir = pagedir_create ();
  if (t->pagedir == NULL)
    goto done;
  process_activate ();

  /* Open executable file. */
  file = filesys_open (file_name);
  ...
}
```

```
void pagedir_activate (uint32_t *pd)
{
  if (pd == NULL)
    pd = init_page_dir;
  asm volatile ("movl %0, %%cr3" : : "r" (vtop (pd)) : "memory");
}
```

```
void process_activate (void)
{
  struct thread *t = thread_current ();
  pagedir_activate (t->pagedir);

  /* Set thread's kernel stack for use in processing interrupts. */
  tss_update ();
}
```

After this point, the user virtual memory mappings changed!

Wait, …

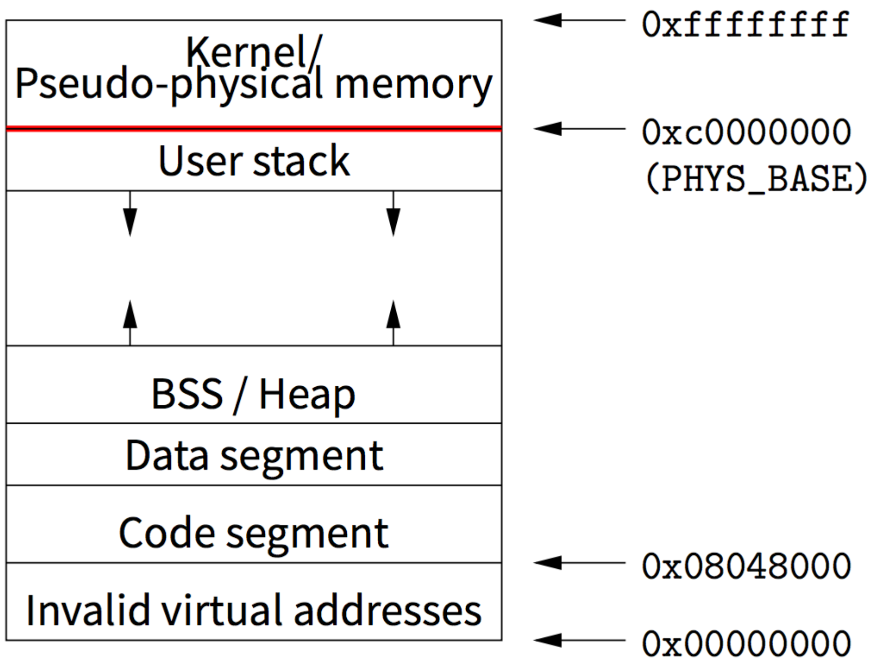
We just changed the user virtual memory mappings, how is it OK for us to still access these variables we created earlier, e.g., `file_name`?

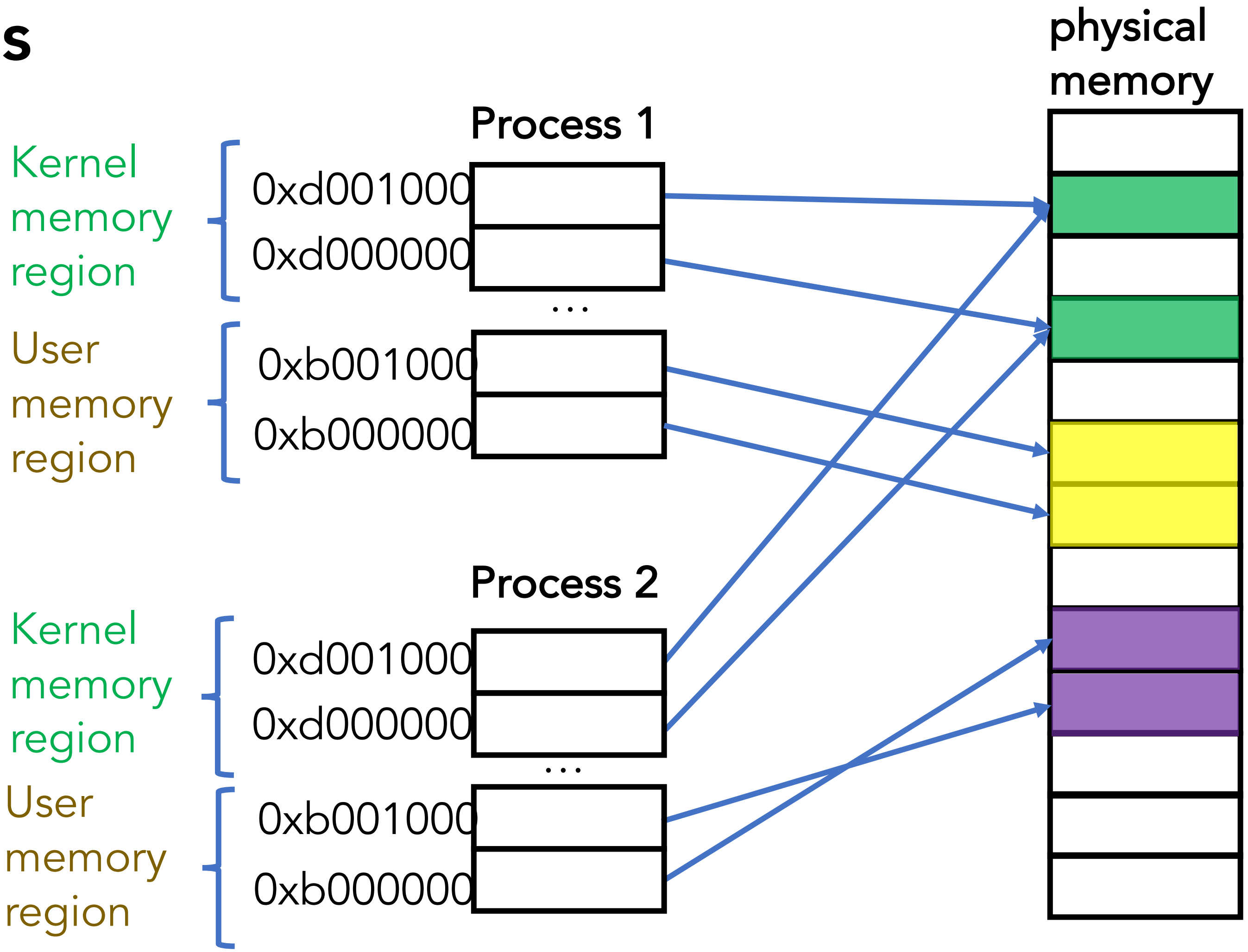
A related concern: how to access variables across multiple processes?

- e.g., to implement `int wait (pid_t pid)` you want to create a variable in `struct thread` to store some information for a process,
  - e.g., thread->wait_status,
- but how can you read/write this variable from the parent process?

Answer: *We're in the Kernel!*

The kernel virtual memory mappings are the same across all processes





Answer: *We're in the Kernel!*

The kernel virtual memory mappings are the same across all processes

Implications:

- When we context switch to another process, although it involves changing the page tables, the kernel virtual memory addresses are still valid after the switch
- All objects created in the kernel functions are accessible across processes
  - e.g.,

```
static struct list all_list;      threadX->wait_status
```

- Memory for user processes will be freed when a user process exits, but memory objects allocated within the kernel code using `malloc` should be explicitly freed!

How Is This Implemented?

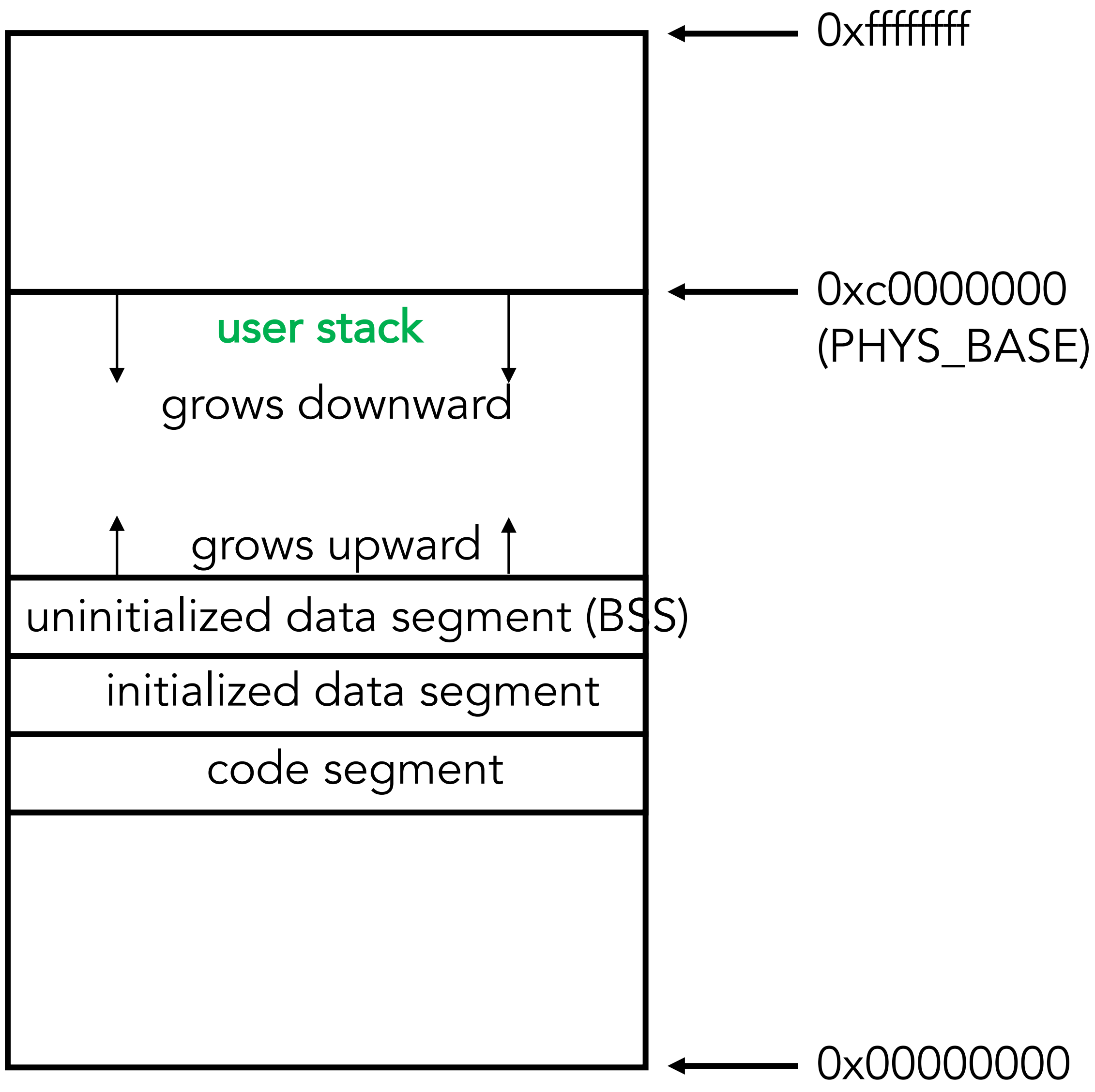
```
bool load (const char *file_name, void (**eip) (void), void **esp)
{
  struct thread *t = thread_current ();
  ...
  /* Allocate and activate page directory. */
  t->pagedir = pagedir_create ();
  if (t->pagedir == NULL)
    goto done;
  process_activate ();

  /* Open executable file. */
  file = filesys_open (file_name);
  ...
}
```

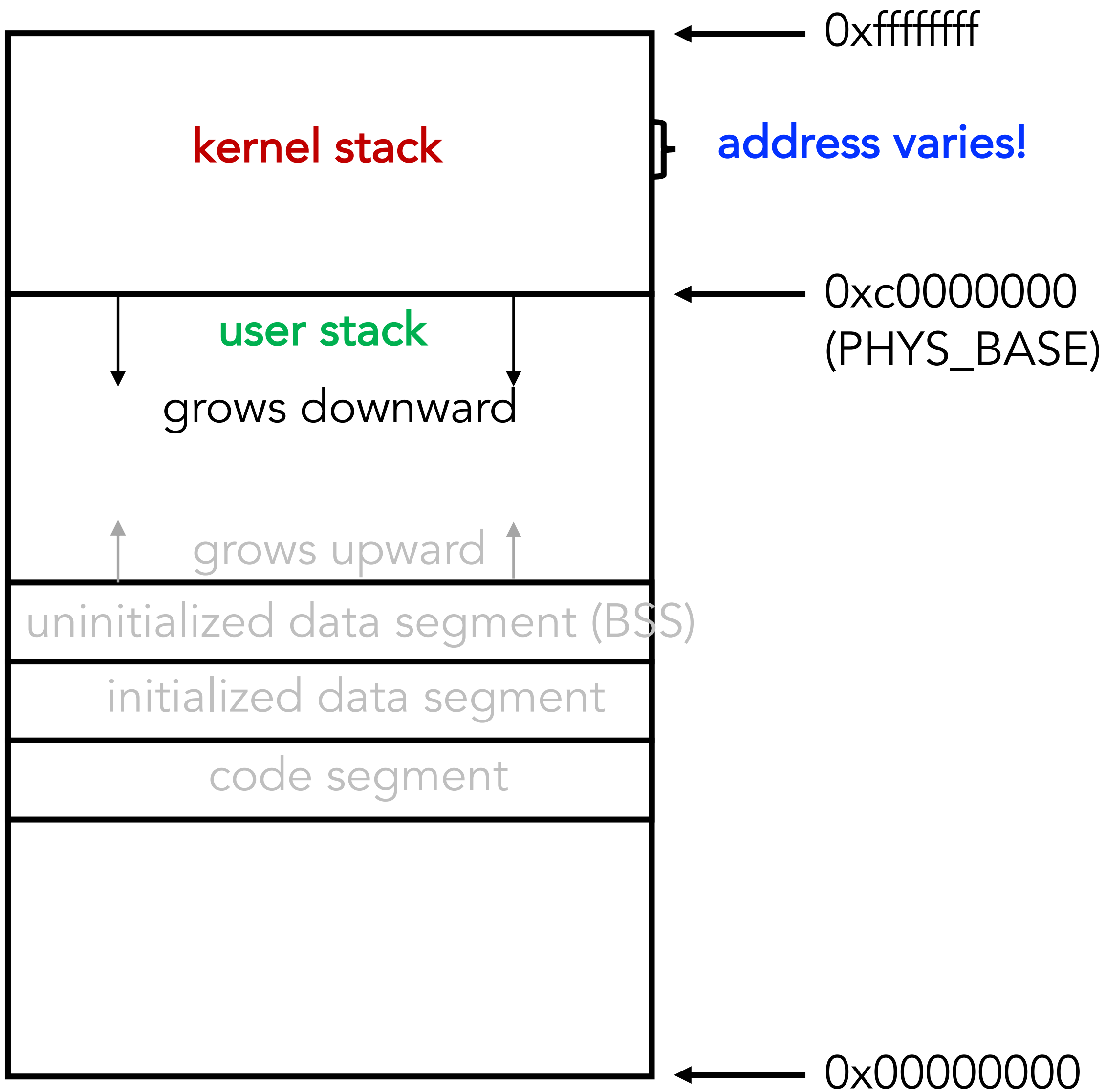
```
uint32_t *
pagedir_create (void)
{
  uint32_t *pd = palloc_get_page (0);
  if (pd != NULL)
    memcpy (pd, init_page_dir, PGSIZE);
  return pd;
}
```

Initialized in [paging_init() in thread.c]

User Stack vs Kernel Stack

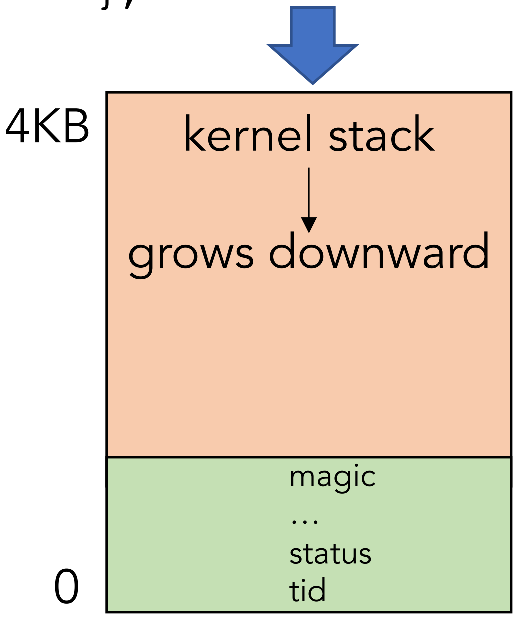


User Stack vs Kernel Stack



A new kernel thread gets a new kernel stack

```
struct thread {
  tid_t tid;
  enum thread_status status;
  …
  unsigned magic;
};
```



Lab 2

Minimal changes to get started:

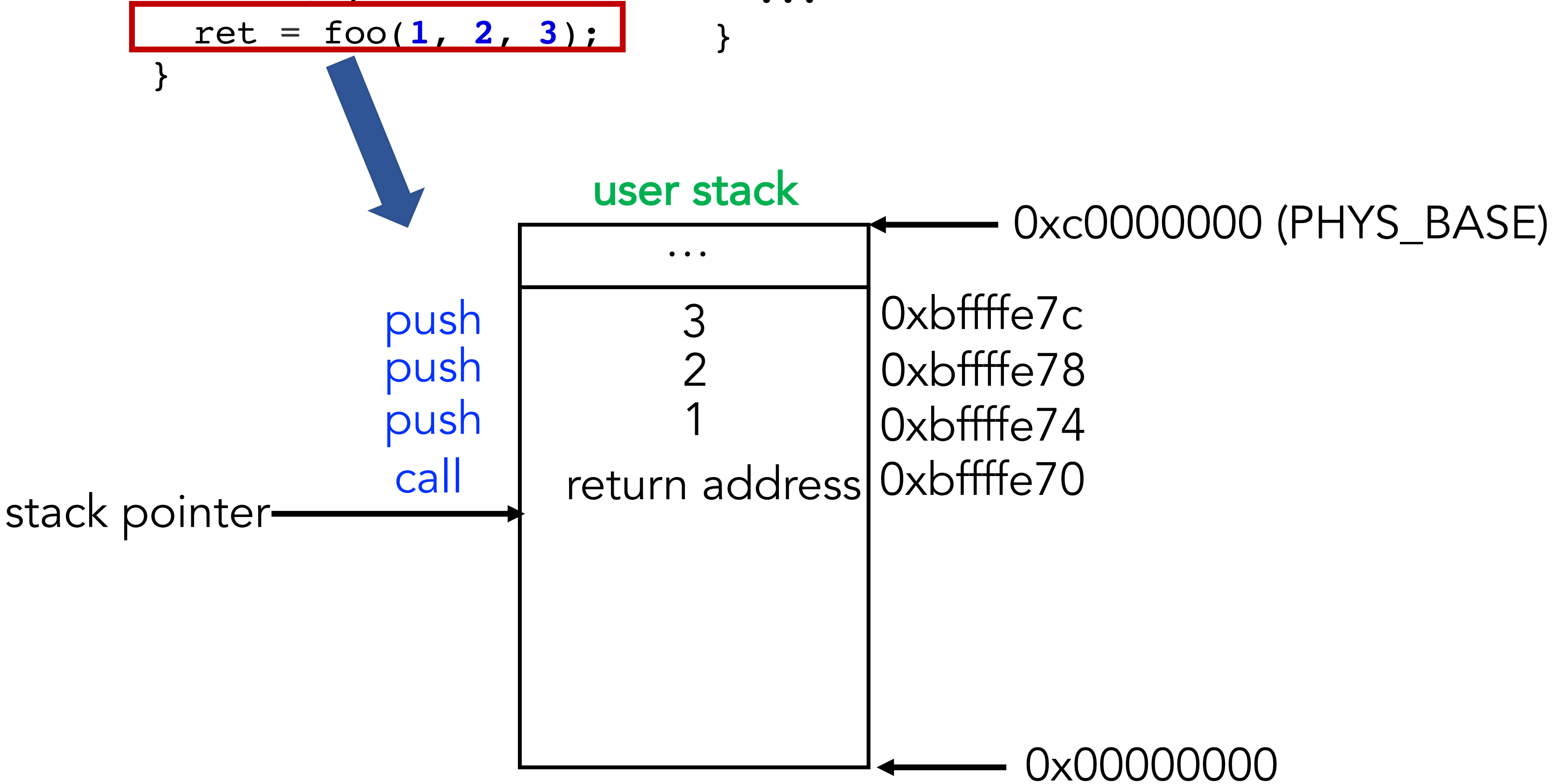
1. setup_stack(): `*esp = PHYS_BASE;` ➔ `*esp = PHYS_BASE - 12;`
2. change process_wait() to an infinite loop

Why setting `esp` to `PHYS_BASE – 12` ?

A temporary setup for obeying x86 calling convention

```
void bar()
{
  int ret;
  ret = foo(1, 2, 3);
}
```

```
int foo(int a, int b, int c)
{
  ...
}
```



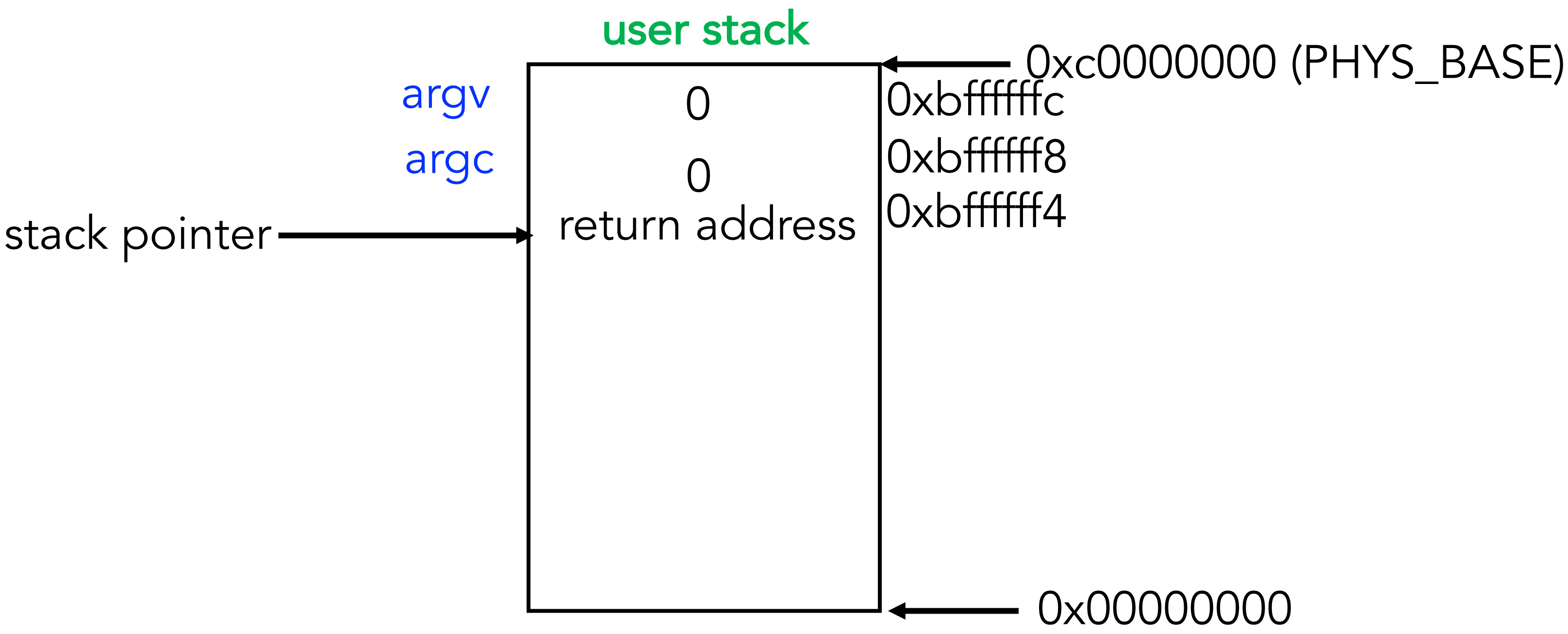
Why setting `esp` to `PHYS_BASE – 12` ?

A temporary setup for obeying x86 calling convention

- Every user program’s entry point is:

```
void _start (int argc, char *argv[])
{
  exit (main (argc, argv));
}
```

- minimal 3 elements on user stack, each 4 bytes = 12



user stack

argv 0 0xbffffffc

argc 0 0xbffffff8

stack pointer → return address 0xbffffff4

0xc0000000 (PHYS_BASE)

0x00000000

Why setting `esp` to `PHYS_BASE – 12` ?

A temporary setup for obeying x86 calling convention

- Every user program’s entry point is:

```
void _start (int argc, char *argv[])
{
  exit (main (argc, argv));
}
```

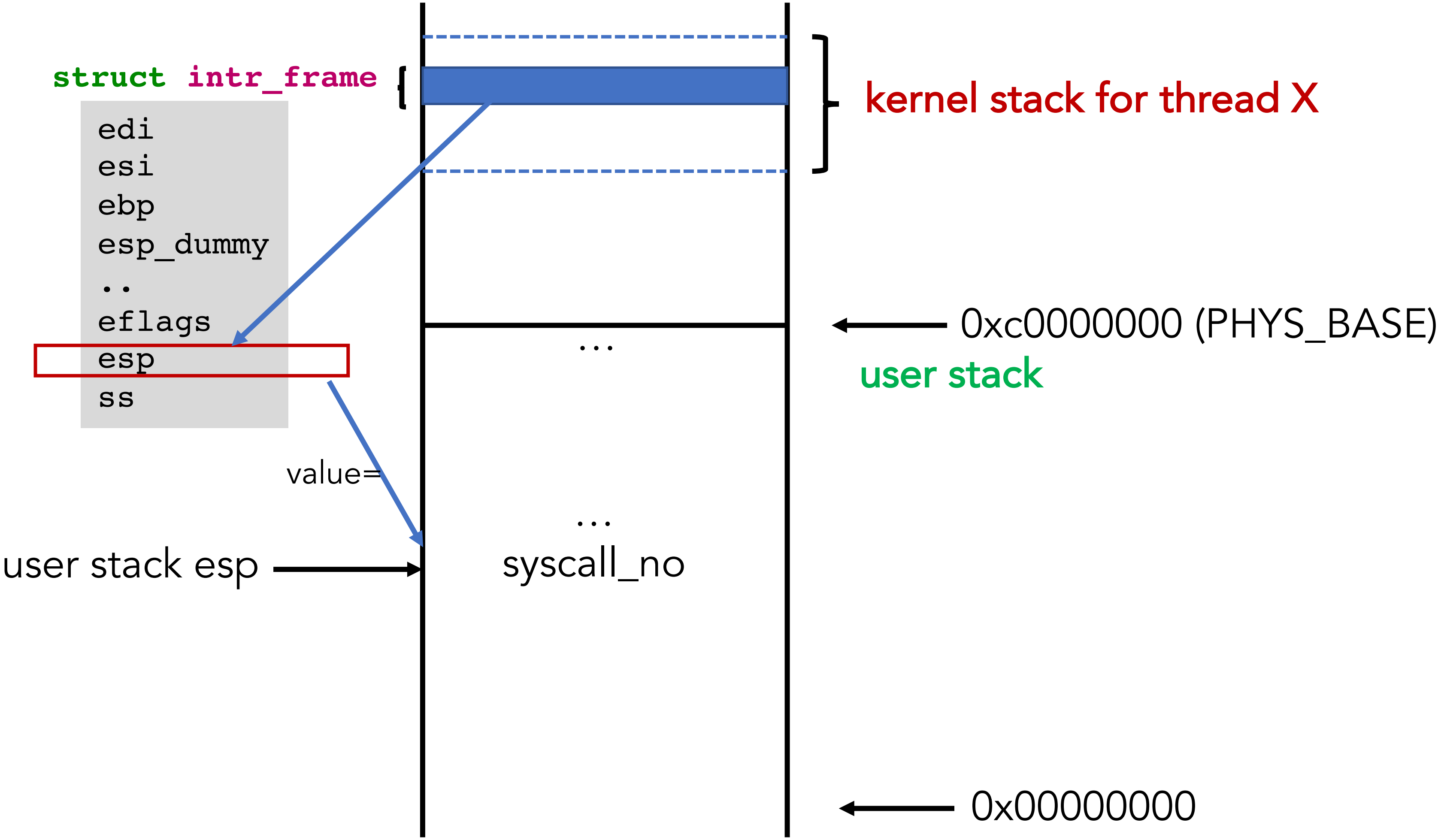
- minimal 3 elements on user stack, each 4 bytes = 12

Note: this is only a temporary setup

- Once you implement argument parsing, you should set esp correctly based on the actual arguments pushed on the user stack

System Call

Through trap (an interrupt frame)



How to retrieve the syscall no in syscall_handler?

from reading user memory at intr_frame->esp

User Memory Access

Upon system call, *no* page directory switch

- i.e., in sycall_handler, the kernel can directly access user memory by dereferencing it
- However, must carefully check each user memory address for robustness!

Two approaches for checking + accessing user memory

- Software approach: using pagedir methods to check validity of an address
  - Easier (straightforward), but slower
- Hardware approach: leveraging page fault to detect invalid address
  - Fast, a bit more difficult to understand (but not difficult to implement once you figure it out)

Hardware Approach

Try loading the memory from a given address `addr`

- Assume addr is the function argument

```
movl 4(%esp), %edx;
movzbl (%edx), %eax;
```

edx = addr
eax = [addr]

- **Problem:** we'll get a page fault if addr is invalid
- **Idea:** let page fault handler inform us, how?

Hardware Approach

Use the given helper function, modify page fault handler

```
/* Reads a byte at user virtual address UADDR.
   UADDR must be below PHYS_BASE.
   Returns the byte value if successful, -1 if a segfault
   occurred. */
static int get_user (const uint8_t *uaddr)
{
  int result;
  asm ("movl $1f, %0; movzbl %1, %0; 1:"
       : "=&a" (result) : "m" (*uaddr));
  return result;
}
```

compilation

```
get_user:
        movl 4(%esp), %edx;
        movl $1f, %eax;
        movzbl (%edx), %eax;
    1:
        ret
```

- If addr is valid, eax has the value
- If addr is invalid, the page fault handler will
  - set eip to address of label 1 (stored in eax now)
  - set eax to be -1 (0xffffffff);
  - resume to ret

Hardware Approach

Use the given helper function, modify page fault handler

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  asm ("movl $1f, %0; movzbl %1, %0; 1:"
       : "=&a" (result) : "m" (*uaddr));
  return result;
}
```

But what if the value at uaddr is -1? We can't tell if it's invalid or not!

- **Solution:** read one byte at a time!
  - If value is valid, at most can be 255 (0xff)
  - How to represent a valid -1? Read four bytes (call get_user four times), convert to an integer!